i-v Fuzzy Shortest Path in a Multigraph

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ABSTRACT

In this research paper the author introduces the notion of i-v fuzzy multigraph. The classical Dijkstra’s algorithmic rule to search out the shortest path in graphs isn’t applicable to fuzzy graph theory. Consequently the author proposes a brand new algorithmic rule referred to as by IVF-Dijkstra’s algorithmic rule with the philosophy of the classical Dijkstra’s formula to unravel the SPP in an i-v fuzzy multigraph

Keywords: i-v fuzzy number, IVFSPA, IVF-Dijkstra’s

INTRODUCTION

The ‘Multigraph’ is a generalized data structure of ‘graph’ where multiple links might exist between a pair of nodes. For example, in a data communication model such as in Adhoc Network or MANET, the multipath or multi route features are very common. Two neighboring routers in a network constellation may share multiple direct connections between them (instead of simply one), therefore as to scale back the bandwidth as compared if one association is employed. Many real world situations of communication network, transportation network, etc. can’t be sculptured into graphs, however are often well sculptured into multigraphs. In many of these directed multigraphs, the weights of the arcs are not always crisp numbers or fuzzy numbers (FNs) but interval-valued fuzzy numbers (i-v fuzzy numbers) depending upon the choice of the concerned decision makers. In this paper, we’ve considered multigraphs without any loops.

One of the first studies on fuzzy shortest path problem (FSPP) in graphs was done by Dubois and Prade and then by Klein. However, few additional solutions to FSPP projected in [16, 36-38] are fascinating, although the work of Dubois and Prade was a serious breakthrough, however that paper lacked any practical interpretation as though fuzzy shortest path is found, however still this might not really be any of the path within the corresponding network for which it had been found, however there’s no
work reported within the existing literature on finding a fuzzy shortest path in a multigraph. In this paper we solve the FSPP problem for a multigraph where the arc-weights are interval valued (i-v) fuzzy numbers. We follow the conception of classical Dijkstra’s algorithmic rule that is applicable to graphs with crisp weights, and then extend this idea to directed multigraphs wherever the weights of the arcs are i-v fuzzy numbers.

**Preliminaries**

For details of the classical notion of fuzzy set theory of Zadeh, one could see any good book on it. The conception of i-v fuzzy numbers is of importance for quantifying associate degree ill-known amount. In our work here throughout, we use the basic operations like fuzzy addition $\oplus$, fuzzy subtraction $\ominus$, and ‘ranking’ of fuzzy numbers, etc. Trivially, any crisp real number can be viewed as a fuzzy number and any fuzzy number can be viewed as an i-v fuzzy number There is no unique method for ranking n number of i-v fuzzy numbers. But there are methods in soft computing for ranking fuzzy numbers. Each technique has some advantages and disadvantages relying upon the properties of the appliance domains and therefore the problem statement into account. However, if $A_1, A_2, A_3, \ldots, A_n$ be n i-v fuzzy numbers ranked in fuzzy ascending order according to any pre determined technique i.e. $A_1 \prec A_2 \prec A_3 \ldots \ldots, \ldots, A_n$ then $A_1$ and $A_n$ are known as severally the i-v fuzzy-min and i-v fuzzy-max of those n i-v fuzzy numbers.

In this section we present basic preliminaries of the theory of multigraphs. A multigraph $X$ is an ordered pair $(V, E)$ which consists of two sets $V$ and $E$, where $V$ or $V(X)$ is the set of vertices (or, nodes), and $E$ or $E(X)$ is the set of edges (or, arcs). Here, though multiple edges or arcs would possibly exist between any two vertices, however we tend to contemplate that no loop exists. Multigraphs are basically of two catagories: undirected and directed. In an undirected multigraph the edge $(i, j)$ and the edge $(j, i)$, if exist, are clearly identical not like within the case of directed multigraph. For a latest algebraic study on the theory of multigraphs, the work and also may be seen. Figure 1 illustrated below a directed multigraph $X = (V, E)$, in which $V = \{G, H, I, J, K\}$ and $E = \{GH_1, GH_2, GI, IG, GJ, IJ, JH, JK, IK, HK, KH\}$.

Any multigraph $Y = (W, F)$ is referred as a submultigraph of the multigraph $X = (V, E)$ iff $W \subseteq V$ and $F \subseteq E$. The Figure 2 illustrated below is a submultigraph $H$ of the given multigraph $G$ shown in Figure 1.

In many of the real world issues of networks, be it during a communication model or transportation model, the weights of the arcs aren’t invariably crisp but fuzzy numbers. for instance, the Figure-3 below shows a public road transportation model for a somebody wherever the value parameter for traveling every arc are obtainable to him as an i-v fuzzy number.

![Fig. 1: A directed multigraph X](image1)

![Fig. 2: A submultigraph H of the directed multigraph G](image2)
In this paper we consider such type of real situations in directed multigraphs of communication systems or transportation systems and develop a method to find an i-v fuzzy shortest path from a source node to a destination node.

i-v Fuzzy Shortest Path in a Directed Multigraph

The FSPP has been solved by various authors for graphs, but for a directed multigraph there is no attempt made so far in the literature for searching a fuzzy shortest path. In our technique here, we tend to solve this FSPP for multigraphs wherever we tend to conjointly use the notion of Dijkstra’s algorithmic rule however with simple soft-computations. We call our proposed algorithm by the name IVF-Dijkstra’s Algorithm.

For developing the IVF-Dijkstra’s Algorithm, first of all we need to define the terms: i-v Fuzzy-Min-Weight arc-set, i-v Fuzzy shortest path estimate \( d(v) \) of a vertex, i-v Fuzzy relaxation of an arc, etc. in the context of the multigraphs, and develop few subalgorithms.

i-v Fuzzy-Min Weight Arc-set of a Directed Multigraph

Consider a directed multigraph \( G \) where the arcs are of i-v fuzzy weights. Suppose that there are \( n \) number of arcs from the vertex \( u \) to the vertex \( v \) in \( G \), where \( n \) is a non-negative integer. Let \( W_{uv} \) be the set whose elements are the arcs between vertex \( u \) and vertex \( v \), but keyed & sorted in non-descending order by the value of their respective i-v fuzzy weights, using a suitable pre-decided ranking method of i-v fuzzy numbers.

\[
w_{uv} = \{ (u_1, w_{1uv}), (u_2, w_{2uv}), \ldots, (u_n, w_{nuv}) \}
\]

where, \( u \) is the arc-\( i \) from vertex \( u \) to vertex \( v \) and \( w \) is the i-v fuzzy weight of it, for \( i = 1, 2, 3, \ldots, n \). If two or more number of i-v fuzzy weights are equal then they may appear at random at the corresponding place of non-descending array with no loss of generality in our discussion.

If there is no confusion, let us denote the multiset \( \{ w_{1uv}, w_{2uv}, \ldots, w_{nuv} \} \) also by the same name \( W_{uv} \). Let \( w_{uv} \) be the i-v fuzzy-min value of the multiset \( W_{uv} = \{ w_{1uv}, w_{2uv}, \ldots, w_{nuv} \} \). Clearly, \( w_{uv} = w_{1uv} \), as the multiset \( W_{uv} \) is already sorted.

Then the set \( W = \{ < (u,v), w_{uv} > : (u,v) \in E \} \) is called the i-v fuzzy-min-weight arc-set of the multigraph \( G \). Suppose that the subalgorithm FMWA \( (G) \) returns the i-v fuzzy-min-weight arc-set \( W \).

i-v fuzzy Shortest path estimate \( d(v) \) of a vertex \( v \) in a directed multigraph

Suppose that node \( s \) is that the source vertex and the presently traversed vertex is node \( u \). There is no single value of weight for arc between vertex \( u \) and a neighbor vertex \( v \), rather there are multiple value of weights as there are multiple arcs between vertex \( u \) and vertex \( v \). Using the value of \( w_{uv} \) from the i-v fuzzy-min weight multiset \( W \) of a directed multigraph, we can now find the i-v fuzzy shortest path estimate i.e. \( d(v) \) of any vertex \( v \), in a directed multigraph as below :-

\[
(i-v \text{ fuzzy shortest path estimate of vertex } v) = (i-v \text{ fuzzy shortest path estimate of vertex } u) \oplus (i-v \text{ fuzzy-min of all the i-v fuzzy weights corresponding to the arcs from the vertex } u \text{ to the vertex } v).
\]

or, \( d(v) = d(u) \cdot w_{uv} \).

i-v fuzzy Relaxation of an Arc

We extend the classical notion of relaxation to the case of i-v fuzzy weighted arcs. By i-v fuzzy relaxation we shall mean here the relaxation process of an arc whose weight is an i-v fuzzy number. For this, first of all we initialize the multigraph along with its starting vertex and i-v fuzzy shortest path estimate for each vertices of the multigraph \( G \). The following ‘I-V FUZZY-INITIALIZATION-SINGLE-SOURCE’ algorithm will do :

I-V FUZZY-INITIALIZATION-SINGLE-SOURCE \( (G, s) \)

1. For each vertex \( v \in V[G] \)
2. \( d[v] = 221Efsa \)
3. \( v. \pi = \text{NIL} \)
4. \( d[s] = 0 \)

After the i-v fuzzy initialization, the process of i-v fuzzy relaxation of each arc begins. The
concerned sub-algorithm I-V FUZZY-RELAX, plays the very important role to update $d[v]$, i.e. the i-v fuzzy shortest traversed cost between the starting node $s$ and the node $v$ (which is the neighbor of the presently traversed vertex $u$).

**I-V FUZZY-RELAX** ($u$, $v$, $W$)
1. IF $d[v] > d[u] \oplus w_{uv}$
2. THEN $d[v] \leftarrow d[u] + w_{uv}$
3. $v.p \leftarrow u$

where, $w_{uv}$ is the i-v fuzzy-min weight of the arcs from vertex $u$ to vertex $v$, and $v.p$ denotes the parent node of vertex $v$.

**i-v fuzzy Shortest Path Algorithm (IVFSPA)**

In this section we now present our main algorithm to find single source i-v fuzzy shortest path in a directed multigraph. We name this I-v fuzzy Shortest Path Algorithm by the title IVFSPA. In this algorithmic rule we tend to use the above subalgorithms, and additionally the subalgorithm EXTRACT-I-V FUZZY-MIN($Q$) that extracts the node $u$ with minimum key by exploiting i-v fuzzy ranking technique and updates $Q$.

**IVFSPA** ($G$, $s$)
1. I-V FUZZY-INITIALIZATION-SINGLE-SOURCE ($G$, $s$)
2. $W \leftarrow$ FMWA ($G$)
3. $S \leftarrow \emptyset$
4. $Q \neq V[G]$
5. WHILE $Q \neq \emptyset$
6. $u \leftarrow$ EXTRACT-I-V FUZZY-MIN ($Q$)
7. $S \leftarrow S \cup \{u\}$
8. FOR each vertex $v \in$ Adj[$u$]
9. $u$-$v$ FUZZY-RELAX ($u$, $v$, $W$)

**Example**

Consider the following directed Multigraph $G$ with i-v fuzzy weights of its arcs. We want to resolve the single-source i-v fuzzy shortest paths problem taking the vertex $A$ as the source node and the vertex $D$ as the destination node.

Clearly the IVFSPA algorithm computes to yield the following result:

1. $w_{AB} = 12^{-}$, $w_{AC} = 6^{-}$, $w_{CB} = 8^{-}$, $w_{CD} = 3^{+}$, and $w_{BD} = 3^{-}$; and then
2. $S = \{A, C, B, D\}$, i.e. the i-v fuzzy shortest
path from the source vertex A is:
A → C → B → D.
3. d-values i.e. i-v fuzzy shortest distance estimate-values of each vertex from the starting vertex A upto the destination vertex D is:

CONCLUSION

Multigraph is an important generalization of the mathematical model 'graph'. There are many real life problems of network, transportation, communication, circuit systems, etc. which cannot be modeled into graphs but into multigraphs only21,32. In many of these directed multigraphs, the weights of the arcs are not always crisp or ordinary fuzzy but i-v fuzzy33,35. The FSPP has been solved for graph by several authors, but not for multigraphs. In this paper we consider the FSPP and have developed a method to find i-v fuzzy shortest path from a source vertex to a destination vertex of a directed multigraph. It is expected that our proposed method and the algorithms for FSPP on multigraphs can play a vital role in many application areas of computer science, communication network, transportation systems, etc. in particular in those networks that can't be sculptured into graphs however into multigraphs.

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370


