Z-Dijkstra’s Algorithm to solve Shortest Path Problem in a Z-Graph

SIDHARTH SANKAR BISWAS*

Department of Computer Science & Engineering Faculty of Engineering & Technology
Jamia Hamdard University Hamdard Nagar, New Delhi – 110062, India.
*Corresponding author E-mail:ssbiswas1984@gmail.com

http://dx.doi.org/10.13005/ojcst/10.01.24

(Received: March 07, 2017; Accepted: March 17, 2017)

ABSTRACT

In this paper the author introduces the notion of Z-weighted graph or Z-graph in Graph Theory, considers the Shortest Path Problem (SPP) in a Z-graph. The classical Dijkstra’s algorithm to find the shortest path in graphs is not applicable to Z-graphs. Consequently the author proposes a new algorithm called by Z-Dijkstra’s Algorithm with the philosophy of the classical Dijkstra’s Algorithm to solve the SPP in a Z-graph.

Keywords: Z-number, Z-distance, Z-weighted graph, Z-graph, Z-Dijkstra’s.

INTRODUCTION

In graph theory, the Shortest Path Problem (SPP) is one of the most famous problems studied and being studied by researchers. The SPP is the problem of finding a path between two vertices (or nodes) in a digraph such that the sum of the weights of its constituent edges is minimized. Thus the core problem is to find the shortest path from a source vertex S to a single destination vertex D in a directed graph and to compute the corresponding min cost. Shortest Path problems (SPP) are among the fundamental problems studied in Computational Geometry, Graph Algorithms, Geographical Information Systems (GIS), Network Optimization etc. to list a few only out of many. Sometimes the network of a real life communication or transportation system can not be modeled into a graph but into a multigraph. In such a case the standard algorithms of graph theory cannot be applied in SPP too. Biswas et al. has done rigorous analysis of SPP in multigraphs. However, in this work we consider those networks which can be modeled into graphs.

The concept of crisp graphs is extended to define fuzzy graphs for soft computing using graph theoretic algorithms. The Fuzzy Shortest Path Problem (FSPP) is a generalization of the classical SPP for applications in ill-defined environment and has been found important to many applications such as Communication or Transportation Network, Computational Geometry, Graph Algorithms, Geographical Information Systems (GIS), Network Optimization etc. to list a few only out of many.
Systems (GIS), Network Optimization, etc. In traditional shortest path problems, the arc length of the network takes precise numbers, but in the real-world problem, the arc length may represent transportation time or cost which can be known only approximately due to vagueness of information, and hence it can be considered a fuzzy number or an intuitionistic fuzzy number. The nodes are well precise, but the data about the weights (costs) of the links are sometimes not available as crisp numbers, rather fuzzy numbers or intuitionistic fuzzy numbers or Z-numbers. Recently Zadeh has introduced the notion of Z-number in 2011, a new direction in soft-computing number theory. In our work here we use the theory of Z-numbers of Zadeh.

One of the first studies on fuzzy shortest path problem (FSPP) in graphs was done by Dubois and Prade and then by Klein. Though the work of Dubois and Prade was a major break-through, but that paper lacked any practical interpretation as even if fuzzy shortest path is found, but still this may not actually be any of the path in the corresponding network for which it was found. According to their approach, the shortest path length can be obtained, but the corresponding path in the network may not exist. This drawback of their method made it difficult for application in network problems. Klein proposed a dynamic programming recursion-based fuzzy algorithm. Lin and Chen found the fuzzy shortest path length in a network by means of a fuzzy linear programming approach. Okada and Soper proposed a fuzzy algorithm, which was based on multiple-labelling methods to offer non-dominated paths to a decision maker. Chuang and Kung proposed a fuzzy shortest path length procedure that can find a fuzzy shortest path length among all possible paths in a network. Yao and Lin presented two new types of fuzzy shortest path network problems. The main results obtained from their studies were that the shortest path in the fuzzy sense corresponds to the actual paths in the network, and the fuzzy shortest path problem is an extension of the crisp case. Nayeem and Pal have proposed an algorithm based on the acceptability index introduced by Sengupta and Pal, which gives a single fuzzy shortest path or a guideline for choosing the best fuzzy shortest path according to the decision maker’s viewpoint. Thus, numerous papers have been published in different journals/books on the fuzzy shortest path problem (FSPP).

In this work the authors introduces Z-graph and solves the SPP in a Z-graph by developing a new algorithm in the style of the famous Dijkstra’s Algorithm to make it applicable in much wider domains.

**Preliminaries of Z-numbers**

The Z-number is a new fuzzy-theoretic concept introduced by Prof. Zadeh in 2011. There are not much work reported so far in the literature about the theory of Z-numbers because of its recent birth and baby stage. But owing to its strong modeling it will surely take a huge role to the scientists and engineers of all the fields in the world.

In this section we reproduce this new concept of Zadeh from It extends the basic philosophy of Computing With Words (CWW) to include the perception of uncertainty of the information conveyed by a natural language statement. The Z-number thus, serves as a model of linguistic summarization of natural language statements, a technique to merge human-affective perspectives with CWW, and consequently can be envisaged to play a radical role in the domain of CWW-based system design and Natural Language Processing (NLP).

**Z-number**

Decisions are based on information. To be useful, information must be reliable. Basically, the concept of a Z-number relates to the issue of reliability of information.

A Z-number Z is an ordered pair, having two components and hence can be expressed using the expression $Z = (A, B)$. The first component, $A$, is a restriction (generalized constraint) on the values which a real-valued uncertain variable, $X$, is allowed to take. The second component, $B$, is a measure of reliability (certainty) of the first component. Typically, $A$ and $B$ are described in a natural language.

Thus a Z-number can be expressed as an ordered pair of fuzzy numbers denoted as $Z = (A, B)$. For simplicity, in our work here $A$ and $B$ are assumed as triangular fuzzy numbers.
Clearly Z-numbers can be used to model uncertain information in real world. For example, in risk analysis, the loss of severity of the fifth component is very low, with a confidence of very likely, which can be written as a Z-number as follows $Z = (\text{very low}; \text{very likely})$. Example of another type of Z-numbers are:

$Z_1 = (\text{about 45 minutes}, \text{very sure}), Z_2 = (\text{about 30 minutes}, \text{sure})$.

A Z-number $Z = (A, B)$ is called to be a null Z-number denoted by $0_z$, if both A and B are null fuzzy numbers.

**Some Operation on Z-numbers**

Consider two Z-numbers $Z_1$ and $Z_2$ given by

$Z_1 = (A_1, B_1)$

and

$Z_2 = (A_2, B_2)$

where $A_i$ and $B_i$ are triangular fuzzy numbers.

There are a number of methods (viz. $^1,^2,^{21},^{29}$) for ranking triangular fuzzy numbers. We follow the Centroid Method because of its simplicity and appropriateness in the philosophy of soft computing.

We say that $Z_1$ is strongly greater than $Z_2$ if $A_1 > A_2$ and $B_1 > B_2$. Otherwise, we say that $Z_1$ is weakly greater than $Z_2$ if $A_1 > A_2$.

Addition $\oplus$ of two Z-numbers $Z_1$ and $Z_2$ yields another Z-number $Z_3$ given by

$Z_3 = Z_1 \oplus Z_2 = (A_1 + A_2, B_1 + B_2)$

As mentioned by Zadeh, it is fact that computation with Z-numbers is an important generalization of computation with real numbers. In particular, the generality of Z-numbers opens the door to construction of better models of reality, especially in fields such as decision analysis, planning, risk assessment, economics and biomedicine.

**Z-weighted graph or Z-graph**

In this section we define Z-graph. A Z-graph is basically a generalized concept of fuzzy graph. Consider the following graph (see Figure 1) where at least one of the weights is Z-number. This type of graph is called a Z-weighted graph or Z-graph (in short).

In this Z-graph in Figure 1, the edge AB has the Z-weight the Z-number (Approx. 67 km, sure), the edge AC has the Z-weight the Z-number (Approx. 49 km, very sure), etc. The z-number weight of an edge is also called the Z-distance between the corresponding two nodes.

A Z-number $Z = (A, B)$ is called to be a null Z-number denoted by $0_z$, if both A and B are null fuzzy numbers.

We say that $Z_1$ is strongly greater than $Z_2$ if $A_1 > A_2$ and $B_1 > B_2$. Otherwise, we say that $Z_1$ is weakly greater than $Z_2$ if $A_1 > A_2$.

Addition $\oplus$ of two Z-numbers $Z_1$ and $Z_2$ yields another Z-number $Z_3$ given by

$Z_3 = Z_1 \oplus Z_2 = (A_1 + A_2, B_1 + B_2)$

As mentioned by Zadeh, it is fact that computation with Z-numbers is an important generalization of computation with real numbers. In particular, the generality of Z-numbers opens the door to construction of better models of reality, especially in fields such as decision analysis, planning, risk assessment, economics and biomedicine.

**Z-weighted graph or Z-graph**

In this section we define Z-graph. A Z-graph is basically a generalized concept of fuzzy graph. Consider the following graph (see Figure 1) where at least one of the weights is Z-number. This type of graph is called a Z-weighted graph or Z-graph (in short).

In this Z-graph in Figure 1, the edge AB has the Z-weight the Z-number (Approx. 67 km, sure), the edge AC has the Z-weight the Z-number (Approx. 49 km, very sure), etc. The z-number weight of an edge is also called the Z-distance between the corresponding two nodes.

![Fig. 1: A transportation network as a Z-graph G](image1)

![Fig. 2: $d[v]$ in a Z-graph](image2)
Z-Dijkstra’s Algorithm for SPP in a Z-graph

The Classical Dijkstra’s algorithm\textsuperscript{27} solves the single-source shortest path problems in a simple graph. In this section we develop a new algorithm called by Z-Dijkstra’s Algorithm (of the style of the classical famous Dijkstra’s algorithm) to solve a SPP in a Z-graph.

Shortest path estimate of a vertex in a directed Z-graph

Consider a Z-weighted directed graph \( G = (V, E) \). Z-shortest path estimate \( d[v] \) of any vertex \( v \), where vertex \( v \) is one of the neighboring vertices of the currently traversed vertex \( u \), is the Z-distance between the vertex \( v \) and vertex \( u \), added with the shortest Z-distance between the starting vertex \( s \) and vertex \( u \), where \( s, u, v \in V[G] \).

\[
\therefore d[v] = (\text{shortest Z-distance between } s \text{ and } u) \oplus (\text{Z-weight of arc between } v \text{ and } u)
\]

This is shown in Figure 2.

Relaxation of an arc in our proposed Z-Dijkstra’s Algorithm

For the relaxation process of an arc to happen, we must first initialize the graph along with its starting vertex \( s \) and Z-shortest path estimate for each vertices of the graph \( G \).

\textbf{INITIALIZE-SINGLE-SOURCE}(G, s)

1. FOR each vertex \( v \in V[G] \)
2. \( d[v] = \infty \)
3. \( v.\pi = NIL \)
4. \( d[s] = 0z \)

(Note : We store all predecessor nodes of \( u \) in the attribute \( u.\pi \). Thus \( s.\pi \) is always Nil, because \( s \) is the source node). Now on the basis of this initialization process, Z-Dijkstra’s algorithm proceeds further and the process of relaxation of each arc begins. The sub-algorithm \( \text{RELAX} \), plays the vital role to update \( d[v] \) i.e. the Z-shortest distance value between the starting vertex \( s \) and the vertex \( v \) (\textit{which is neighbor of the current traversed vertex } \( u, \forall u, v \in V[G] \))

The RELAX algorithm runs as shown below :

\textbf{RELAX} (u, v, w)

1. IF \( d[v] > d[u] \oplus w(u,v) \)
2. THEN \( d[v] \leftarrow d[u] \oplus w(u,v) \)
3. \( v.\pi \leftarrow u \)

where, \( w(u,v) \) is the Z-weight of the arc from vertex \( u \) and vertex \( v \), and \( v.\pi \) denotes the parent node of a vertex \( v \), \( \forall v \in V[G] \). This is shown below in the Figure 3.

Z-Dijkstra’s algorithm solves the single-source shortest-path on a Z-weighted directed Z-graph \( G = (V, E) \) for the case in which all edge Z-weights are non-negative. Z-Dijkstra’s algorithm maintains a set \( S \) of vertices whose final Z-shortest path weights from the source vertex \( s \) has already been determined. The algorithm repeatedly selects the vertex \( u \in V \setminus S \) with the minimum Z-shortest-path estimate, adds \( u \) to \( S \), and relaxes all edges leaving \( u \). Our proposed algorithm is as follows:

\textbf{Z-DIJKSTRA} (G, w, s)

1. \textbf{INITIALIZE-SINGLE-SOURCE} (G, s)
2. \( S \leftarrow \emptyset \)
3. \( Q \leftarrow V[G] \)
4. WHILE \( Q \neq \emptyset \)
5. \( u \leftarrow \text{EXTRACT-MIN}(Q) \)
6. \( S \leftarrow S \cup \{u\} \)
7. FOR each vertex \( v \in \text{Adj}[u] \)
8. \( \text{DO RELAX}(u, v, w) \)

\begin{figure}[ht]
\centering
\includegraphics[width=\textwidth]{fig3.png}
\caption{Diagram showing how RELAX algorithm works.}
\end{figure}
CONCLUSION

Z-graph is a generalization of fuzzy graph. There are many real life problems of networks in communication systems, transportation, circuit systems, etc. which cannot be modeled into graphs but into fuzzy graphs or Z-graphs only. The classical Dijkstra's algorithm to find the shortest path in graphs is not applicable to Z-graphs. In this paper we have done slight adjustment in the classical famous Dijkstra's algorithm to make it applicable to Z-graphs to find the Z-shortest path from a source vertex to a destination vertex. The modified algorithm is called as Z-Dijkstra's algorithm. The networks where the classical Dijkstra's algorithm fails can now be dealt with the Z-Dijkstra's algorithm for performing efficient and optimal communication/transportation in many cases.

REFERENCES

1. Abbasbandy, Saeid, Ranking of fuzzy numbers, some recent and new formulas, IFSA-EUSFLAT, 2009, Page 642-646.


37. Mahdavi, I., Tajdin, A., Nourifar, R.,


